GD 2023 Weakly and Strongly Fan-Planar Graphs



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Fan-Planarity

Each edge *e* can only be crossed by a *fan* of edges, i.e., a bundle of edges sharing a common endpoint that all cross *e* from the same side.

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Anchor of e

























[Cheong et al., 2022]





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Do the graph classes of weak & strong fan-planarity coincide?





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Adjacency-Crossing

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[Brandenburg, 2020] Adjacency-crossing graphs with *m* edges: construction of fan-planar graph with the same # of edges

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First step, we eliminate certain configurations:

Input: Let G = (V, E) be a (preprocessed) graph on *n* vertices;

admitting a weakly fan-planar drawing Γ ; with the least number of pattern (III)

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Induction step (roughly): G is split into two weakly fan-planar graphs G_1 and G_2 , amounting combined to the same number of edges as G, which contain less triples forming pattern (III):





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Weak Fan-Planarity: Edge Density Induction on the # of edge triples forming pattern (III) G_2 $|E(G)| = |E(G_1)| + |E(G_2)|$ G_1 e_{l} e \mathcal{U}





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Happy to take your questions



Thank you!