

# Minimizing the energy of spherical graph representations

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Graph Drawing 2023

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# Graph Representations

## Definition

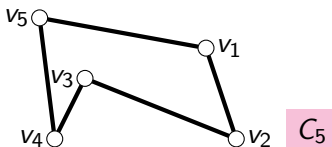
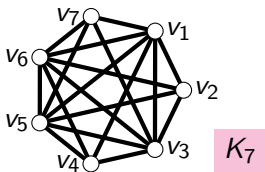
A *representation of a graph  $G$  in  $\mathbb{R}^d$*  is a function  $\mathbf{r} : V(G) \rightarrow \mathbb{R}^d$ .  
The point  $\mathbf{r}(v)$  is the *representation of  $v \in V(G)$* .



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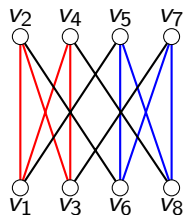
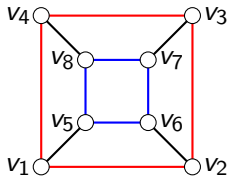
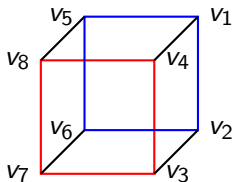
## Definition

A *representation of a graph  $G$  in  $\mathbb{R}^d$*  is a function  $r : V(G) \rightarrow \mathbb{R}^d$ . The point  $r(v)$  is the *representation of  $v \in V(G)$* .



A representation induces a straight line/geometric drawing.

# Graph Representations: Symmetry, Planarity, Partition



# Graph Representations in $\mathbb{R}^d$

The *energy* of a representation  $\mathbf{r} : V(G) \rightarrow \mathbb{R}^d$  of a graph  $G$  is defined as

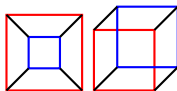
$$\text{energy}(G, \mathbf{r}) = \sum_{uv \in E(G)} \|\mathbf{r}(u) - \mathbf{r}(v)\|^2.$$

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- energy **minimization** was used for **plane graph drawings** (Tutte) and for **symmetrical representations** (spectral method)

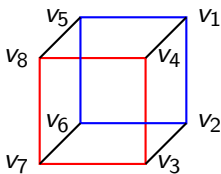


- energy **maximization** was used for **graph partitions** (Goemans-Williamson)



## Definition

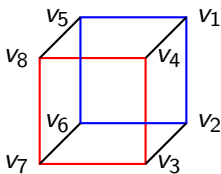
A *representation matrix* of a representation  $\mathbf{r}$  of  $G$  in  $\mathbb{R}^d$  is a  $d \times v(G)$  matrix where column  $v \in V(G)$  is  $\mathbf{r}(v)$ .



$$V = \begin{bmatrix} \mathbf{r}(v_1) & \mathbf{r}(v_2) & \mathbf{r}(v_3) & \mathbf{r}(v_4) & \mathbf{r}(v_5) & \mathbf{r}(v_6) & \mathbf{r}(v_7) & \mathbf{r}(v_8) \end{bmatrix}$$

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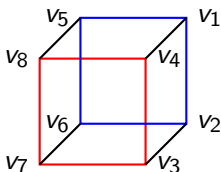


$$V = \begin{bmatrix} 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 \end{bmatrix}$$



Minimize the energy w.r.t

the rows of the representation matrix are orthogonal, unit vectors



$$V = \frac{1}{\sqrt{8}} \begin{bmatrix} 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 \end{bmatrix}$$

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**Algorithm 1:** Spectral Graph Drawing Algorithm.

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**Input:** Adjacency matrix  $A$

- 1 Compute an orthogonal, unit eigenvector basis for the  $d + 1$  highest eigenvalues of  $A$ .
  - 2 Let  $R$  be the representation matrix where the rows are the basis vectors.
  - 3  $v \in V(G)$  is represented by column  $v$  of  $R$ .
- 

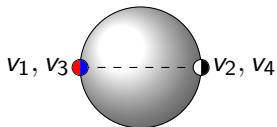
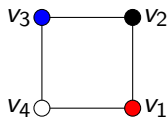


Drawings of platonic solids use  $\lambda_2$ -eigenspace.

Goemans, Williamson, 1995

1. Maximize the energy w.r.t

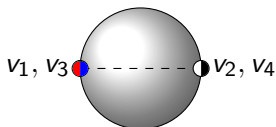
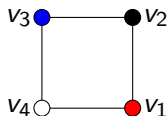
the vertices are on the unit  $d$ -sphere for some  $d \geq 1$ .



## Goemans, Williamson, 1995

## 1. Maximize the energy w.r.t

the vertices are on the **unit  $d$ -sphere** for some  $d \geq 1$ .



## 2. Partition vertices by taking a random hyperplane of the sphere.

**Theorem (Goemans, Williamson, 1995)**

*The algorithm (1.+2.) produces an edge-cut with expected size at least 0.87856 times the size of the maximum edge cut of the graph.*

# Minimizing on the sphere

DeVos, Rogers, W., 2023

Minimize the energy w.r.t

the vertices are on a **unit  $d$ -sphere** for some  $d \geq 1$  (1)

the origin is the barycenter of the vertex representations (2)

This minimization can be expressed via a **semidefinite program** (and hence approximated efficiently).



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A *unit barycentre*  $\mathbf{0}$  representation is a representation which satisfies (1)+(2).

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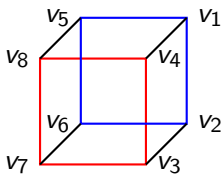
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$$V = \frac{1}{\sqrt{3}} \begin{bmatrix} 1 & 1 & -1 & -1 & 1 & 1 & -1 & -1 \\ 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & -1 & -1 & 1 & 1 & -1 & -1 & 1 \end{bmatrix}$$

# Random Regular Graphs

## Theorem (DeVos, Rogers, W., 2023)

For every random  $d$ -regular graph  $G$  there exists a *minimum energy, unit barycentre  $\mathbf{0}$  representation  $\mathbf{r}$*  such that the inequalities

$$(d - 2\sqrt{d-1} - \epsilon)v(G) \leq \text{energy}(G, \mathbf{r}) \leq (d - 2\sqrt{d-1} + \epsilon)v(G).$$

hold asymptotically almost surely for every  $\epsilon > 0$ .



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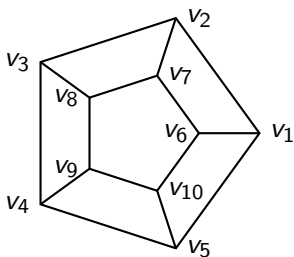
This means

$$\min_{r: \text{unit, barycentre } \mathbf{0}} \text{energy}(G, \mathbf{r}) \approx (d - \lambda_2)v(G).$$

## Vertex-transitive Graphs

If for any  $u, v \in V(G)$  there exists an automorphism that maps  $u$  to  $v$ , then  $G$  is called *vertex-transitive*.

An *automorphism* of a graph  $G$  is a permutation  $\sigma$  of  $V(G)$ , such that  $(u, v)$  is an edge if and only if  $(\sigma(u), \sigma(v))$  is an edge.



$$\sigma = \begin{pmatrix} v_1 & v_2 & v_3 & v_4 & v_5 & v_6 & v_7 & v_8 & v_9 & v_{10} \\ v_7 & v_8 & v_9 & v_{10} & v_6 & v_2 & v_3 & v_4 & v_5 & v_1 \end{pmatrix}$$

## Theorem (DeVos, Rogers, W., 2023)

*For each connected vertex-transitive graph  $G$  there exists a minimum energy, unit barycentre  $\mathbf{0}$  representation  $\mathbf{r}$  of  $G$*

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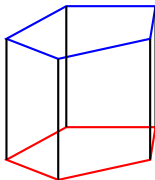
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Construction of representation:

- $w$  is an eigenvector to eigenvalue  $\lambda_2$
- $w_\sigma$  is obtained from  $w$  by permuting entries w.r.t.  $\sigma$
- $R$  is a matrix with rows  $(w_\sigma)_{\sigma \in \text{Aut}(G)}$

# Our Semidefinite Program

Maximize:  $A \bullet (V^T V)$

Subject To:

$C_v \bullet (V^T V) = 1$  for  $v \in V(G)$

$J \bullet (V^T V) = 0$

minimizes energy

columns have norm 1

origin is barycenter

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$J$  has only 1's

Random projection of representation onto  $d = 2$ .

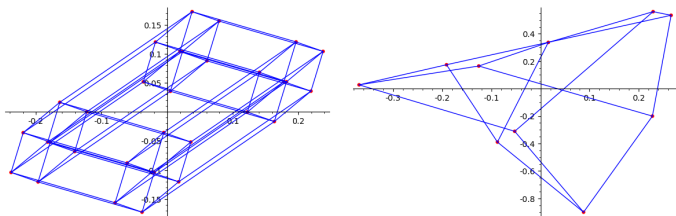


Figure 1: Hypercube for  $d = 4$  and Petersen Graph

Thank you!